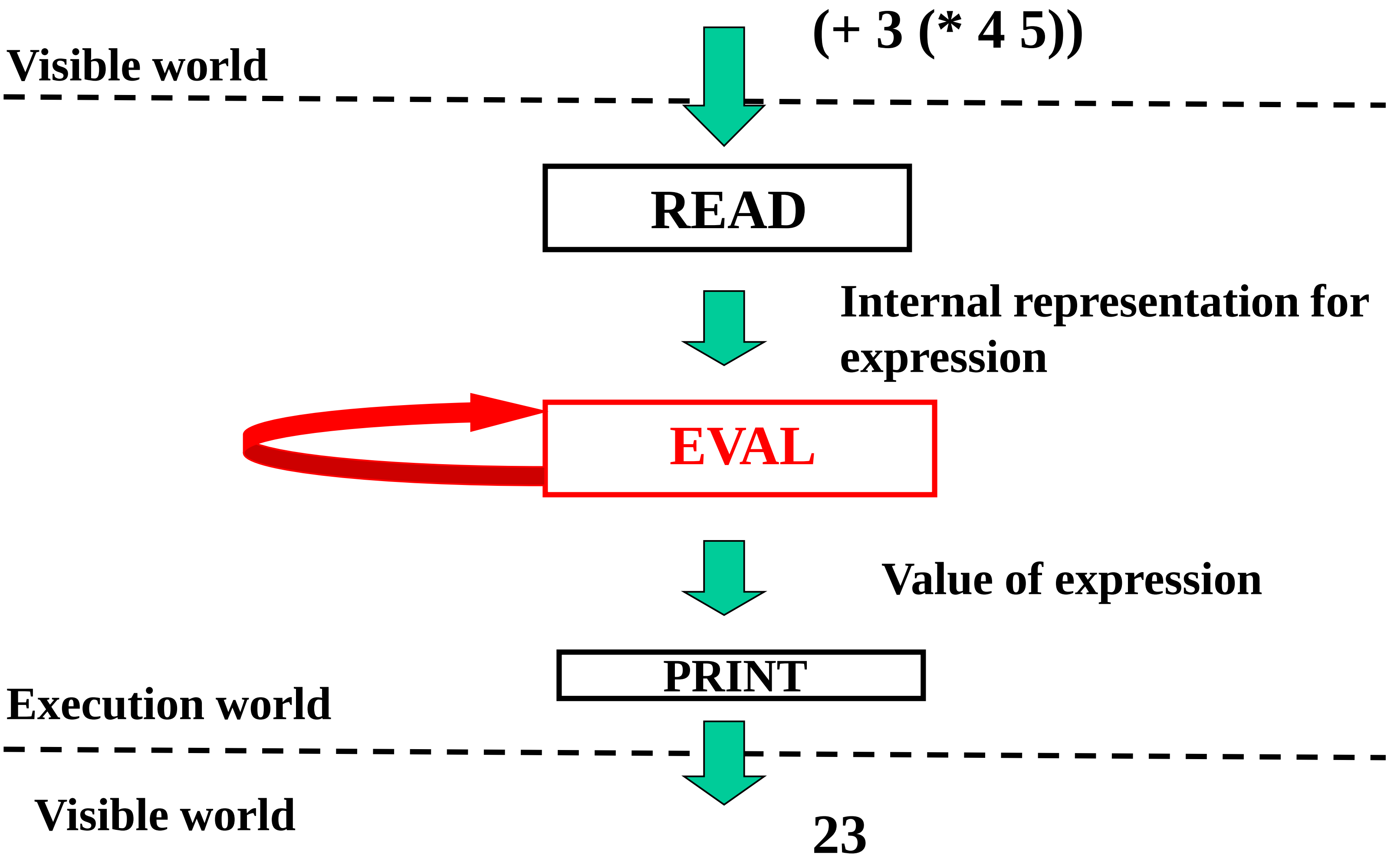
Basic Scheme February 8, 2007

- Compound expressions
- Rules of evaluation
- Creating procedures by capturing common patterns

Previous lecture

- Basics of Scheme
  - Expressions and associated values (or syntax and semantics)
    - Self-evaluating expressions
      - 1, “this is a string”, #f
    - Names
      - +, *, >=, <
    - Combinations
      - (* (+ 1 2) 3)
    - Define
- Rules for evaluation

Read-Eval-Print



Summary of expressions

- Numbers: value is expression itself
- Primitive procedure names: value is pointer to internal hardware to perform operation
- “Define”: has no actual value; is used to create a binding in a table of a name and a value
- Names: value is looked up in table, retrieving binding
- Rules apply recursively

Simple examples

25	→	25
(+ (* 3 5) 4)	→	60
+	→	[#primitive procedure …]
(define foobar (* 3 5))	→	no value, creates binding of foobar and 15
foobar	→	15 (value is looked up)
(define fred +)	→	no value, creates binding
(fred 3 5)	→	15

This lecture

Adding procedures and procedural abstractions to capture processes

Language elements -- procedures

- Need to capture ways of doing things – use procedures

parameters

(lambda (x) (* x x))

body

To process something multiply it by itself

•Special form – creates a procedure and returns it as value

Language elements -- procedures

- Use this anywhere you would use a procedure

((lambda(x)(* x x)) 5)

(* 5 5) **lambda exp** **arg**

25

Language elements -- abstraction

- Use this anywhere you would use a procedure

((lambda(x)(* x x)) 5)

Don’t want to have to write obfuscatory code – so can give the lambda a name

(define square (lambda (x) (* x x)))
(square 5) → 25

Rumplestiltskin effect!
(*The power of naming things*)

Scheme Basics

- Rules for *evaluating*

1. If **self-evaluating,** return value.
2. If a **name,** return value associated with name in environment.
3. If a **special form,** do something special.
4. If a **combination,** then

   a. *Evaluate* all of the subexpressions of combination (in any order)

   b. *apply* the operator to the values of the operands (arguments) and return result

- Rules for *applying*

1. If procedure is **primitive procedure,** just do it.
2. If procedure is a **compound procedure,** then:
   evaluate the body of the procedure with each formal parameter replaced by the corresponding actual argument value.

Interaction of define and lambda

```
1. (lambda (x) (* x x))
            ==> #[compound-procedure 9]

2. (define square (lambda (x) (* x x)))
                                   ==> undef

3. (square 4) ==> 16

4. ((lambda (x) (* x x)) 4) ==> 16

5. (define (square x) (* x x)) ==> undef
```

This is a convenient shorthand (called "syntactic sugar") for 2 above – this is a use of lambda!

Lambda special form

- lambda syntax **`(lambda (x y) (/ (+ x y) 2))`**

- 1st operand position: the parameter list **`(x y)`**
  - a list of names (perhaps empty) **`()`**
  - determines the number of operands required

- 2nd operand position: the body **`(/ (+ x y) 2)`**
  - may be any expression(s)
  - not evaluated when the lambda is evaluated
  - evaluated when the procedure is applied
  - value of body is value of last expression evaluated
- mini-quiz: `(define x (lambda ()(+ 3 2)))`
- `x`
- `(x)`

- semantics of lambda:

THE VALUE OF

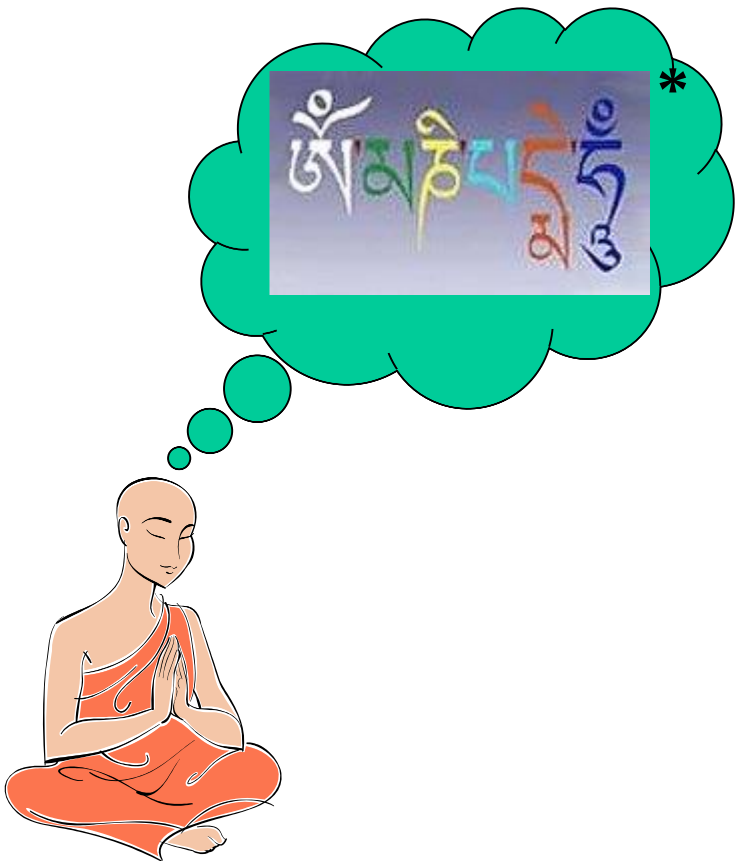
A LAMBDA EXPRESSION

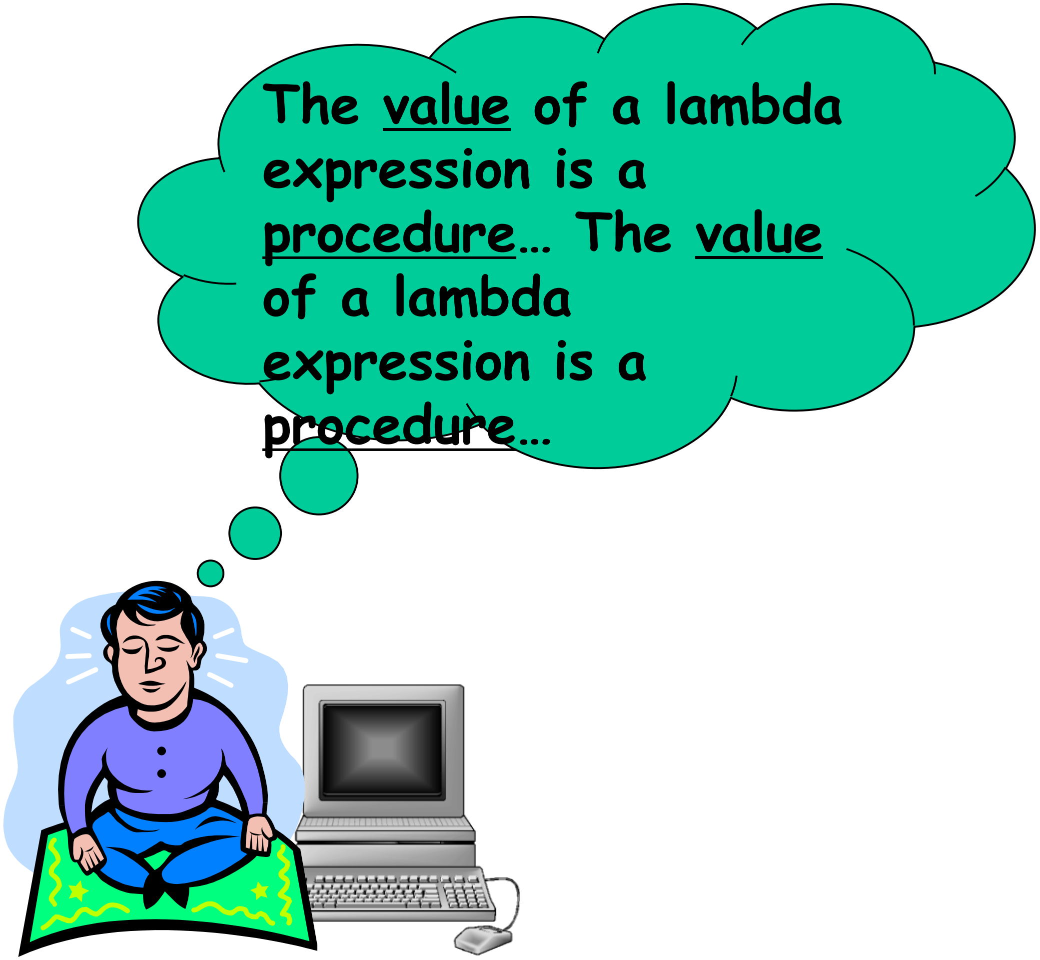
IS

A PROCEDURE

Achieving Inner Peace

(and a good grade)





Om Mani Padme Hum…

Using procedures to describe processes

- How can we use the idea of a procedure to capture a computational process?

What does a procedure describe?

- Capturing a common pattern
  - (* 3 3)
  - (* 25 25)
  - (* foobar foobar)

Common pattern to capture

```
(lambda (x) (* x x) )
```

Name for thing that changes

Modularity of common patterns

Here is a common pattern:

```
(sqrt (+ (* 3 3) (* 4 4)))
(sqrt (+ (* 9 9) (* 16 16)))
(sqrt (+ (* 4 4) (* 4 4)))
```

Here is one way to capture this pattern:

```
(define pythagoras
  (lambda (x y)
    (sqrt (+ (* x x) (* y y)))))
```

Modularity of common patterns

Here is a common pattern:

So here is a cleaner way of capturing the pattern:

```
(define square (lambda (x) (* x x)))
(define pythagoras
  (lambda (x y)
     (sqrt (+ (square x) (square y)))))
```

Why?

- Breaking computation into modules that capture commonality
  - Enables reuse in other places (e.g. square)
- Isolates (abstracts away) details of computation within a procedure from use of the procedure
  - Useful even if used only *once* (i.e., a unique pattern)

```
(define (comp x y)(/(+(* x y) 17)(+(+ x y) 4))))

(define (comp x y)(/ (prod+17 x y) (sum+4 x y)))
```

Why?

- May be many ways to divide up

```
(define square (lambda (x) (* x x)))
(define pythagoras
    (lambda (x y)
        (sqrt (+ (square x) (square y)))))

(define square (lambda (x) (* x x)))
(define sum-squares
    (lambda (x y) (+ (square x) (square y))))
(define pythagoras
    (lambda (y x) (sqrt (sum-squares y x))))
```

Abstracting the process

- Stages in capturing common patterns of computation
  - Identify modules or stages of process
  - Capture each module within a procedural abstraction
  - Construct a procedure to control the interactions between the modules
  - Repeat the process within each module as necessary

A more complex example

- Remember our method for finding sqrts
  - To find the square root of X
    - Make a guess, called G
    - If G is close enough, stop
    - Else make a new guess by averaging G and X/G

The stages of “SQRT”

- When is something “close enough”
- How do we create a new guess
- How do we control the process of using the new guess in place of the old one

Procedural abstractions

For “close enough”:

```
(define close-enuf?
  (lambda (guess x)
    (< (abs (- (square guess) x)) 0.001)))
```

Note use of procedural abstraction!

Procedural abstractions

For “improve”:

```
(define average
    (lambda (a b) (/ (+ a b) 2)))
(define improve
    (lambda (guess x)
        (average guess (/ x guess))))
```

Why this modularity?

- "Average" is something we are likely to want in other computations, so only need to create once
- Abstraction lets us separate implementation details from use
  - Originally:

```
(define average
  (lambda (a b) (/ (+ a b) 2)))
```

  - Could redefine as

```
(define average

  (lambda (x y) (* (+ x y) 0.5)))
```

  - No other changes needed to procedures that use **`average`**
  - Also note that variables (or parameters) are internal to procedure – cannot be referred to by name outside of scope of lambda

Controlling the process

- Basic idea:
  - Given X, G, want **`(improve G X)`** as new guess
  - Need to make a decision – for this need a new *special form*

```
(if <predicate> <consequence> <alternative>)
```

The `IF` special form

```
(if <predicate> <consequence> <alternative>)
```

- Evaluator first evaluates the `<predicate>` expression.
- If it evaluates to a TRUE value, then the evaluator evaluates and returns the value of the `<consequence>` expression.
- Otherwise, it evaluates and returns the value of the `<alternative>` expression.
- **Why must this be a special form? (i.e. why not just a regular lambda procedure?)**

Controlling the process

- Basic idea:
  - Given X, G, want `(improve G X)` as new guess
  - Need to make a decision – for this need a new *special form*

```
(if <predicate> <consequence> <alternative>)
```

  - So heart of process should be:

```
(if (close-enuf? G X)
    G
                   (improve G X)        )
```

  - But somehow we want to use the value returned by "improving" things as the new guess, and repeat the process

Controlling the process

- Basic idea:
  - Given X, G, want `(improve G X)` as new guess
  - Need to make a decision – for this need a new *special form*

```
(if <predicate> <consequence> <alternative>)
```

  - So heart of process should be:

```
(define sqrt-loop (lambda G X)
   (if (close-enuf? G X)
       G
       (sqrt-loop (improve G X) X      )
```

  - But somehow we want to use the value returned by “improving” things as the new guess, and repeat the process
  - Call process `sqrt-loop` and reuse it!

Putting it together

- Then we can create our procedure, by simply starting with some initial guess:

```
(define sqrt
    (lambda (x)
        (sqrt-loop 1.0 x)))
```

Checking that it does the "right thing"

- Next lecture, we will see a formal way of tracing evolution of evaluation process
- For now, just walk through basic steps
  - **`(sqrt 2)`**
    - **`(sqrt-loop 1.0 2)`**
    - **`(if (close-enuf? 1.0 2)`** … …**`)`**
    - **`(sqrt-loop (improve 1.0 2) 2)`**

    This is just like a normal combination
    - **`(sqrt-loop 1.5 2)`**
    - **`(if (close-enuf? 1.5 2)`** … …**`)`**
    - **`(sqrt-loop 1.4166666 2)`**
- And so on…

Abstracting the process

- Stages in capturing common patterns of computation
  - Identify modules or stages of process
  - Capture each module within a procedural abstraction
  - Construct a procedure to control the interactions between the modules
  - Repeat the process within each module as necessary

Summarizing Scheme

- Primitives
  - Numbers **1, -2.5, 3.67e25**
  - Strings
  - Booleans
  - Built in procedures ***, +, -, /, =, >, <,**
  - **-- Names**

Creates a loop in system – allows abstraction of name for object

- Means of Combination
  - (procedure argument$_1$ argument$_2$ … argument$_n$)
- Means of Abstraction
  - Lambda

    Create a procedure
  - Define

    Create names
- Other forms
  - if

    Control order of evaluation